

This is the best assignment solution uploaded with some modification of the following student:

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Question No. 1:

- a) Consider a point-to-point link 2 km in length. At what bandwidth would propagation delay (at a speed of 2×10^8 m/s) equal transmit delay for 100-byte packets? What about 512-byte packets?
- b) How long does it take to transmit x KB over a y-Mbps link? Give your answer as a ratio of x and y.
- c) Suppose that a certain communications protocol involves a per-packet overhead of 100 bytes for headers and framing. We send 1 million bytes of data using this protocol; however, one data byte is corrupted and the entire packet containing it is thus lost. Give the total number of overhead + loss bytes for packet data sizes of 1000, 5000, 10,000, and 20,000 bytes. Which size is optimal?

Solution Part (a):

100-Byte Packets

Given Data:

Distance = 2 Km = 2×10^3 m

SpeedOfLight = 2×10^8 m/s

Packet Size = 100 bytes = $100 \times 8 = 800$ bits

Propagation Delay = Transmit Delay

To Find:

Bandwidth = X =?

Solution:

According to the Formula given in “Computer Networks: A Systems Approach” by “Larry L. Peterson and Bruce S. Davie”;

$$\text{Transmit Delay} = \frac{\text{Size}}{\text{Bandwidth}} \quad \text{and}$$

$$\text{Propagation Delay} = \frac{\text{Distance}}{\text{SpeedOfLight}}$$

Putting the given values in the above equation; we get

$$\text{Transmit Delay} = \frac{800 \text{ b}}{X \text{ b/s}}$$

$$\text{Propagation Delay} = \frac{2 \times 10^3 m}{2 \times 10^8 m/s} = 1 \times 10^{-5} s = 10 \times 10^{-6} s = 10 \mu s$$

$$\text{Propagation Delay} = 10 \mu s$$

Equating the Transmit Delay and the Propagation Delay; we get

$$\text{Propagation Delay} = \text{Transmit Delay}$$

$$\frac{800 b}{X b/s} = 10 \times 10^{-6}$$

$$X \times 10 \times 10^{-6} = 800$$

$$X = \frac{800}{10 \times 10^{-6}} = 80 \times 10^6 \text{ bit/s} = 80 \text{ Mbit/s}$$

$$X = 80 \text{ Mbit/s}$$

Hence, the propagation delay equals transmit delay for 100-byte packets at the bandwidth of 80Mbit/s.

512-Byte Packets

Given Data:

$$\text{Distance} = 2 \text{ Km} = 2 \times 10^3 \text{ m}$$

$$\text{SpeedOfLight} = 2 \times 10^8 \text{ m/s}$$

$$\text{Packet Size} = 512 \text{ bytes} = 512 \times 8 = 4096 \text{ bits}$$

$$\text{Propagation Delay} = \text{Transmit Delay}$$

To Find:

$$\text{Bandwidth} = X = ?$$

Solution:

According to the Formula given in “Computer Networks: A Systems Approach” by “Larry L. Peterson and Bruce S. Davie”;

$$\text{Transmit Delay} = \frac{\text{Size}}{\text{Bandwidth}} \quad \text{and}$$

$$\text{Propagation Delay} = \frac{\text{Distance}}{\text{SpeedOfLight}}$$

Putting the given values in the above equation; we get

$$\text{Transmit Delay} = \frac{4096 b}{X b/s}$$

$$\text{Propagation Delay} = \frac{2 \times 10^3 m}{2 \times 10^8 m/s} = 1 \times 10^{-5} s = 10 \times 10^{-6} s = 10 \mu s$$

$$\text{Propagation Delay} = 10 \mu s$$

Equating the Transmit Delay and the Propagation Delay; we get

$$\text{Propagation Delay} = \text{Transmit Delay}$$

$$\frac{4096 \text{ b}}{X \text{ b/s}} = 10 \times 10^{-6}$$

$$X \times 10 \times 10^{-6} = 4096$$

$$X = \frac{4096}{10 \times 10^{-6}} = 409.6 \times 10^6 \text{ bit/s} = 409.6 \text{ Mbit/s}$$

$$X = 409.6 \text{ Mbit/s}$$

Hence, the propagation delay equals transmit delay for 512-byte packets at the bandwidth of 409.6 Mb/s.

Solution Part (b):

Number of bytes in a KB = 1024 bytes

Number of bits in a byte = 8 bits

$$x \text{ KB} = 1024 \times 8 \times x \text{ bits}$$

Number of bits in a Mb = 10^6 bits

$$y \text{ Mbps} = 10^6 y \text{ bps}$$

The time “t” taken to transmit x KB over a y-Mbps link can be given as follows:

$$t = \frac{8 \times 1024 \times x}{y \times 10^6}$$

$$t = \frac{8192 \times x}{y \times 10^6}$$

$$t = \frac{8192 \times 10^{-6} \times x}{y}$$

$$t = \frac{8.192 \times 10^{-3} \times x}{y}$$

$$t = \frac{8.192x}{y} \text{ ms}$$

Hence, the taken to transmit x KB over a y-Mbps link is found to be 8.912x/y ms.

Solution Part (c):

Given Data:

Data Size of Packet = D

Per-Packet Overhead = 100 bytes

Number of bytes sent using the protocol = 10^6

$$\text{Number of Packets needed} = N = \frac{10^6}{D}$$

$$\text{Overhead} = 100 \times N$$

$$\text{Loss} = D$$

$$\text{Overhead} + \text{Loss} = 100 \times N + D$$

$$\text{Overhead} + \text{Loss} = 100 \times \left(\frac{10^6}{D} \right) + D$$

To Find:

Total number of overhead + loss bytes for packet data size, D=1000 =?

Total number of overhead + loss bytes for packet data size, D=5000 =?

Total number of overhead + loss bytes for packet data size, D=10,000 =?

Total number of overhead + loss bytes for packet data size, D=20,000 =?

Solution:

Overhead + Loss bytes for D=1000

According to the formula given above:

$$\text{Overhead} + \text{Loss} = 100 \times \left(\frac{10^6}{D} \right) + D$$

Putting the given values, we get;

$$\text{Overhead} + \text{Loss} = 100 \times \left(\frac{10^6}{1000} \right) + 1000$$

$$\text{Overhead} + \text{Loss} = 100 \times 10^3 + 1000$$

$$\text{Overhead} + \text{Loss} = 100 \times 1000 + 1000 = 101000 \Rightarrow \text{Overhead} + \text{Loss} = 101000$$

Overhead + Loss bytes for Packet Data Size of 1000 = 101000

Overhead + Loss bytes for D=5000

According to the formula given above:

$$\text{Overhead} + \text{Loss} = 100 \times \left(\frac{10^6}{D} \right) + D$$

Putting the given values, we get;

$$\text{Overhead} + \text{Loss} = 100 \times \left(\frac{10^6}{5000} \right) + 5000$$

$$\text{Overhead} + \text{Loss} = 100 \times \left(\frac{1000}{5} \right) + 5000$$

$$\text{Overhead} + \text{Loss} = 100 \times 200 + 5000 = 25000 \Rightarrow \text{Overhead} + \text{Loss} = 25000$$

Overhead + Loss bytes for Packet Data Size of 5000 = 25000

Overhead + Loss bytes for D=10,000

According to the formula given above:

$$Overhead + Loss = 100 \times \left(\frac{10^6}{D} \right) + D$$

Putting the given values, we get;

$$Overhead + Loss = 100 \times \left(\frac{10^6}{10000} \right) + 10000$$

$$Overhead + Loss = 10000 + 10000 = 20000 \Rightarrow \text{Overhead} + \text{Loss} = 20,000$$

Overhead + Loss bytes for Packet Data Size of 10,000 = 20,000

Overhead + Loss bytes for D=20,000

According to the formula given above:

$$Overhead + Loss = 100 \times \left(\frac{10^6}{D} \right) + D$$

Putting the given values, we get;

$$Overhead + Loss = 100 \times \left(\frac{10^6}{20000} \right) + 20000$$

$$Overhead + Loss = 100 \times \left(\frac{100}{2} \right) + 20000$$

$$Overhead + Loss = 100 \times 50 + 20000$$

$$Overhead + Loss = 5000 + 20000 = 25000 \Rightarrow \text{Overhead} + \text{Loss} = 25,000$$

Overhead + Loss bytes for Packet Data Size of 20,000 = 25,000

Summarizing the results in a table, we have;

Packet data Size (D)	Overhead + Loss Bytes
1000	101,000
5000	25,000
10,000	20,000 → Optimal Packet data Size
20,000	25,000

Hence, the total number of overhead + loss bytes for packet data sizes of 1000, 5000, 10,000, and 20,000 bytes are found to be 101,000 bytes, 25,000 bytes, 20,000 bytes, 25,000 bytes respectively. Since, the packet data size of 10,000 has the lowest number of Overhead + Loss Bytes i.e. 20,000 bytes, it is the optimal packet data size.

Question No. 2:

- a) Calculate the latency (from first bit sent to last bit received) for the following:
- 10-Mbps Ethernet with a single store-and-forward switch in the path, and a packet size of 5000 bits. Assume that each link introduces propagation delay of $10\ \mu\text{s}$ and that the switch begins retransmitting immediately after it has finished receiving the packet.
 - Same as (a) but with three switches.
- b) The 10 Mbps Ethernet segment can be as long as 2500 meters. If the 10 Mbps Ethernet frame specifications stay the same, how does an increase in supported bandwidth affect the segment length?

Solution Part (a):**(a) Latency With One Switch****Given Data:**

Bandwidth = 10-Mbps
 Packet Size = 5000 bits
 Propagation Delay = $10\ \mu\text{s}$
 Number of switches = 1
 Number of links = 2

To Find:

Latency = ?

Solution:

The transmission Delay can be calculated as follows:

$$\text{Transmit Delay} = \frac{\text{Packet Size}}{\text{Bandwidth}}$$

Putting the given values, we get;

$$\text{Transmit Delay} = \frac{5000}{10 \times 10^6} = \frac{5}{10 \times 10^3} = \frac{1}{2 \times 10^3} = 0.0005 = 500 \times 10^{-6} = 500\ \mu\text{s}$$

The Latency can be calculated from the formula as follows:

Latency = Transmit Delay for each link + Propagation Delay for each link

Putting the given values of Propagation and Transmit delay, we get;

Latency = $500\ \mu\text{s}$ for each link + $10\ \mu\text{s}$ for each link

$$\text{Latency} = 500\ \mu\text{s} \times 2 + 10\ \mu\text{s} \times 2 = 1000\ \mu\text{s} + 20\ \mu\text{s} = 1020\ \mu\text{s}$$

Latency = $1020\ \mu\text{s}$

Hence, latency (from first bit sent to last bit received) is found to be $1020\ \mu\text{s}$

(b) Latency With Three Switches**Given Data:**

Bandwidth = 10-Mbps

Packet Size = 5000 bits

Propagation Delay = 10 μ s

Number of switches = 3

Number of links = 4

To Find:

Latency =?

Solution:

The transmission Delay can be calculated as follows:

$$\text{Transmit Delay} = \frac{\text{Packet Size}}{\text{Bandwidth}}$$

Putting the given values, we get;

$$\text{Transmit Delay} = \frac{5000}{10 \times 10^6} = \frac{5}{10 \times 10^3} = \frac{1}{2 \times 10^3} = 0.0005 = 500 \times 10^{-6} = 500 \mu\text{s}$$

The Latency can be calculated from the formula as follows:

Latency = Transmit Delay for each link + Propagation Delay for each link

Putting the given values of Propagation and Transmit delay, we get;

Latency = 500 μ s for each link + 10 μ s for each link

$$\text{Latency} = 500 \mu\text{s} \times 4 + 10 \mu\text{s} \times 4 = 2000 \mu\text{s} + 40 \mu\text{s} = 2040 \mu\text{s} = 2.04 \times 10^{-3} = 2.04 \text{ms}$$

$$\text{Latency} = 2.040 \text{ms}$$

Hence, latency (from first bit sent to last bit received) is found to be 2.04 ms.

Solution Part (b):

Two factors affect the segment length while increasing the supported bandwidth, i.e. the “minimum frame length” and the “contention period”. Increasing the supported bandwidth will result in less transmission time of a 64-byte frame however; the “contention period” should remain the same. By keeping the “contention period” the same, and if the propagation time of a bit stays the same, the length of the segment drops down. Hence, increase in supported bandwidth will decrease the segment length.

Question No. 3:

Read Beej's guide to Network Programming and modify the talker/listener (talker.c and listener.c) according to the following specifications (Listener and talker are communicating using UDP socket):

Listener:

1. The listener should run in continuous mode, i.e., it should handle one talker after another (Make listener iterative server)
2. The listener will bind on port # 4500
3. The listener should maintain a log named "TalkerMsgLog.txt" of all the talkers and their messages. The log should contain the following information
 - a. IP address of the talker
 - b. Port number of the talker
 - c. Message sent by the talker
4. When a listener receives a message from a talker, it save the talkers message in the log and also sends a reply to the talker containing the following well message:
"Welcome: TALKER-IP-ADDRESS, TALKER-PORTNO"
5. When a listener receives a message from a talker, it saves the talker's message in the log "FirstDataMsg: TALKER-IP-ADDRESS, TALKER-PORTNO: this is first formal data message from Talker" and also sends a reply to the talker containing the following message:
"IDMsg: Your First Data Message is received, please send your name and ID No"
6. When the listener receives a message from the talker, it saves the talker's message in the log "IDMsgReply: TALKER-IP-ADDRESS, TALKER-PORTNO: I am Muhammad – MS090400001" and also sends a reply to the talker containing the following message:
"TermMsg: Your Name and ID is logged on Listener, communication terminated"
7. Listener then waits for some new talker and does not terminate.

Talker:

1. A talker takes the following command line argument from the user:
 - a. IP address of the listener
 - b. Port No. of Listener
2. The command line arguments entered by the user will be : ./talker 127.0.0.1 4500
3. Create a log file "ListernerMsgLog.txt"
4. Talker formats a message and sends to the listener on a IP address and port taken as a command line argument.
5. When the talker receives the welcome message as a reply from the listener, it displays the message on the screen and writes it in the log file.
6. Talker formats a message and sends to the listener: I am "Your Name – MS**04000**"(ID).
7. When the talker receives the TermMsg message from the listener, it displays the message on the screen and writes in the log file, terminates the communication and exit.

Solution:

Listener.C

```

/*
** listener.c -- a datagram sockets "server" demo
*/
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <errno.h>
#include <string.h>
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
#include <arpa/inet.h>
#define MYPORT 4500 // the port users will be connecting to
#define MAXBUFLEN 100
int main(void)
{
    int sockfd;
    struct sockaddr_in my_addr; // my address information
    struct sockaddr_in their_addr; // connector's address information
    int addr_len, numbytes;
    char buf[MAXBUFLEN];
    if ((sockfd = socket(AF_INET, SOCK_DGRAM, 0)) == -1) {
        perror("socket");
        exit(1);
    }
    my_addr.sin_family = AF_INET; // host byte order
    my_addr.sin_port = htons(MYPORT); // short, network byte order
    my_addr.sin_addr.s_addr = INADDR_ANY; // automatically fill with my IP
    memset(&(my_addr.sin_zero), '\0', 8); // zero the rest of the struct
    if (bind(sockfd, (struct sockaddr *)&my_addr,
        sizeof(struct sockaddr)) == -1) {
        perror("bind");
        exit(1);
    }
    addr_len = sizeof(struct sockaddr);
    if ((numbytes=recvfrom(sockfd, buf, MAXBUFLEN-1, 0,
        (struct sockaddr *)&their_addr, &addr_len)) == -1) {
        perror("recvfrom");
        exit(1);
    }
}

```

```

}
printf("got packet from %s\n",inet_ntoa(their_addr.sin_addr));
printf("packet is %d bytes long\n",numbytes);
buf[numbytes] = '\0';
printf("packet contains \"%s\"\n",buf);
close(sockfd);
return 0;
}

```

Talker.C

```

/*
** talker.c -- a datagram "client" demo
*/

#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <errno.h>
#include <string.h>
#include <sys/types.h>
#include <sys/socket.h>
#include <netinet/in.h>
#include <arpa/inet.h>
#include <netdb.h>

#define SERVERPORT "4950"    // the port users will be connecting to

int main(int argc, char *argv[])
{
    int sockfd;
    struct addrinfo hints, *servinfo, *p;
    int rv;
    int numbytes;

    if (argc != 3) {
        fprintf(stderr,"usage: talker hostname message\n");
        exit(1);
    }

    memset(&hints, 0, sizeof hints);
    hints.ai_family = AF_UNSPEC;
    hints.ai_socktype = SOCK_DGRAM;

    if ((rv = getaddrinfo(argv[1], SERVERPORT, &hints, &servinfo)) != 0) {

```

```

        fprintf(stderr, "getaddrinfo: %s\n", gai_strerror(rv));
        return 1;
    }

    // loop through all the results and make a socket
    for(p = servinfo; p != NULL; p = p->ai_next) {
        if ((sockfd = socket(p->ai_family, p->ai_socktype,
                            p->ai_protocol)) == -1) {
            perror("talker: socket");
            continue;
        }

        break;
    }

    if (p == NULL) {
        fprintf(stderr, "talker: failed to bind socket\n");
        return 2;
    }

    if ((numbytes = sendto(sockfd, argv[2], strlen(argv[2]), 0,
                          p->ai_addr, p->ai_addrlen)) == -1) {
        perror("talker: sendto");
        exit(1);
    }

    freeaddrinfo(servinfo);

    printf("talker: sent %d bytes to %s\n", numbytes, argv[1]);
    close(sockfd);

    return 0;
}

```

Question No. 4:

Read the attached research paper entitled “Design and Development of a UDP- Based Connection-Oriented Multi-Stream One-to-Many Communication Protocol” with this assignment very carefully. Summarize the research paper into your own words (450-500) in the following headings:

1. Proposed technique
2. Advantages / Benefits
3. Disadvantages / Drawbacks

4. Suggestions & Recommendations

Solution:

Please! Consult the mentioned Research Paper.

Question No. 5:

Read the attached research paper entitled “A Survey of Packet Loss Recovery Techniques for Streaming Audio” with this assignment very carefully, and answer the following questions:

1. Classify the recovery techniques discussed in this paper and write a brief note.
2. Give the comparison of the Forward Error Correction Techniques to repair losses of data during transmission, discussed in this paper.
3. Give the comparison of the Error Concealment Techniques, discussed in this paper.

Solution:

Please! Consult the mentioned Research Paper.

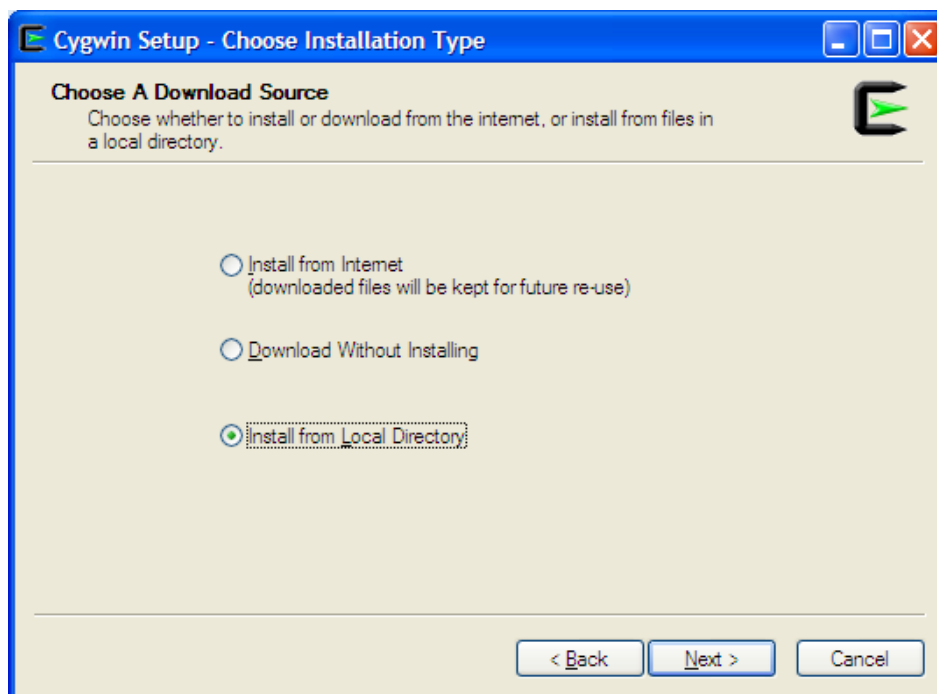
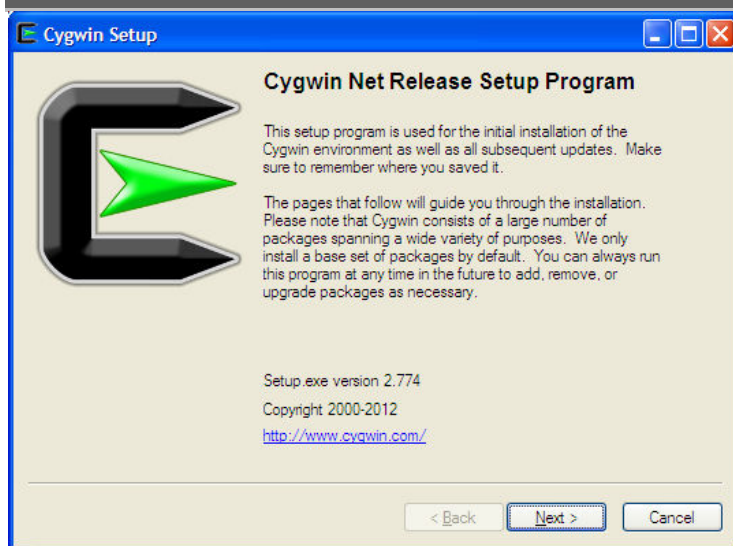
Question No. 6:

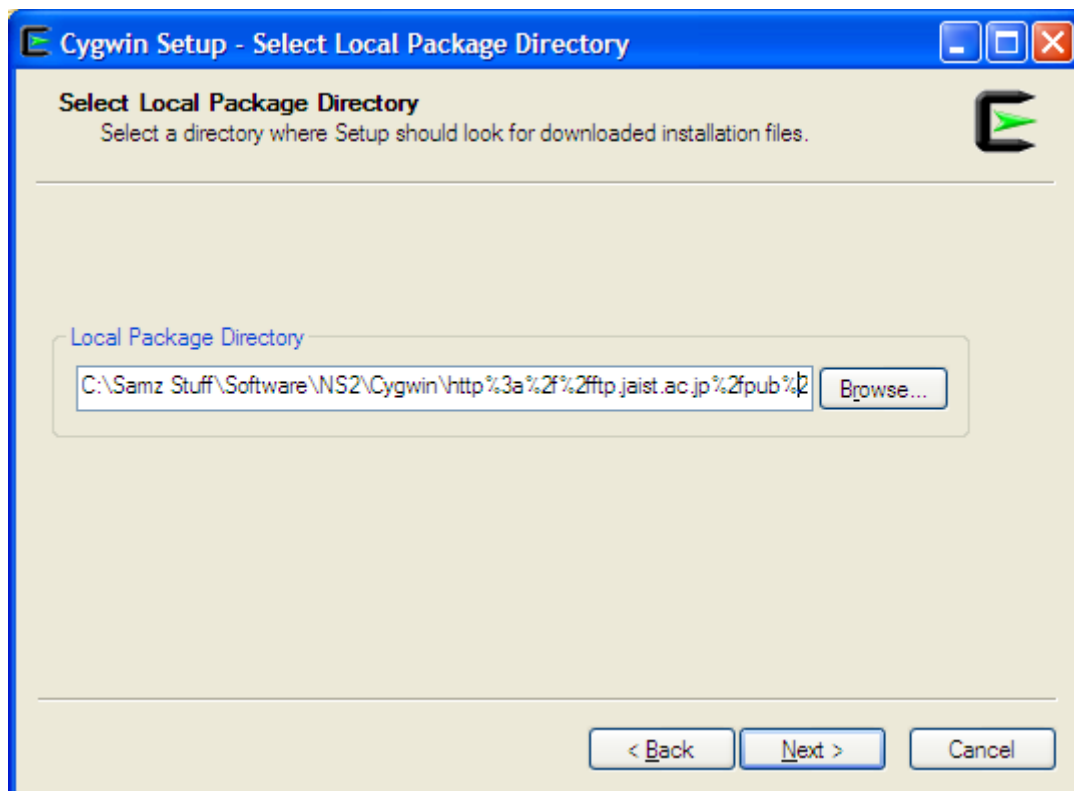
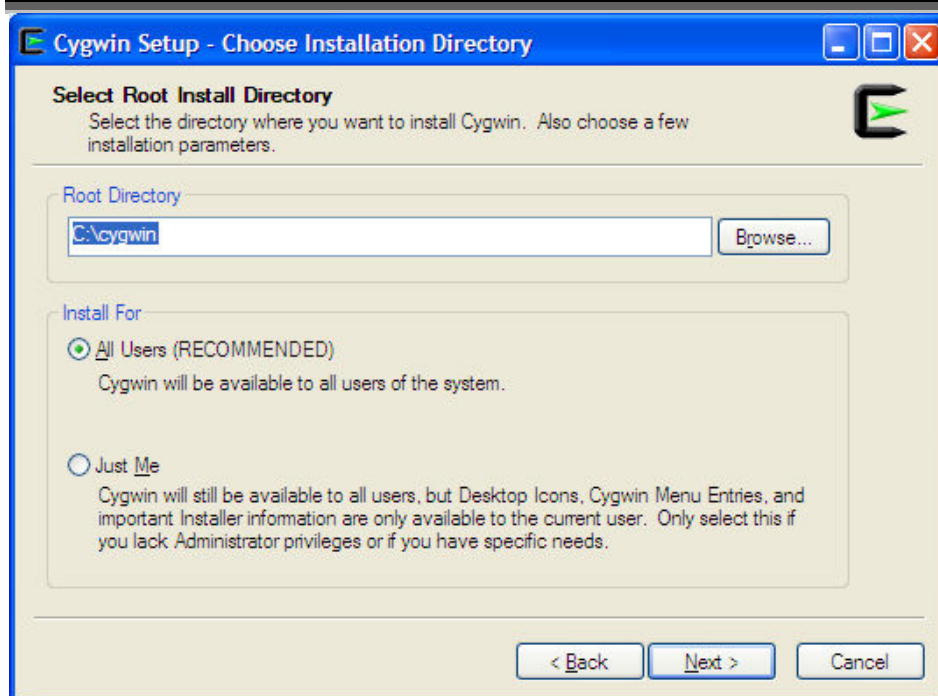
Explore Network Simulator (NS 2).

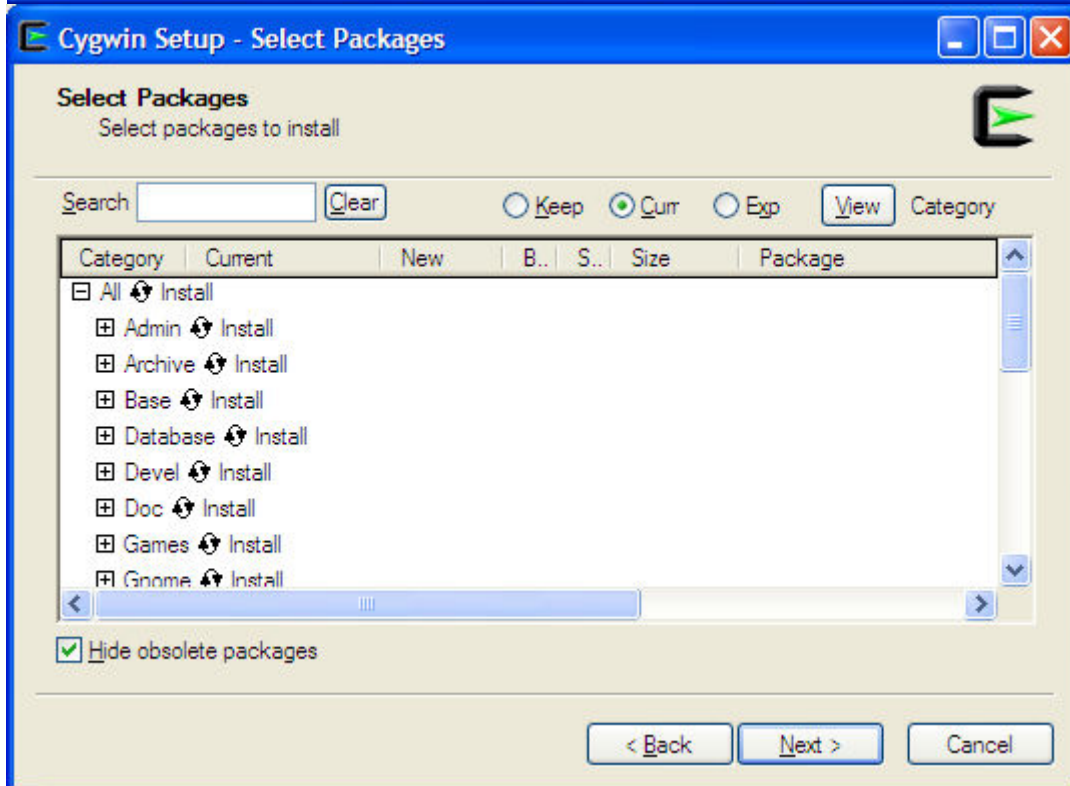
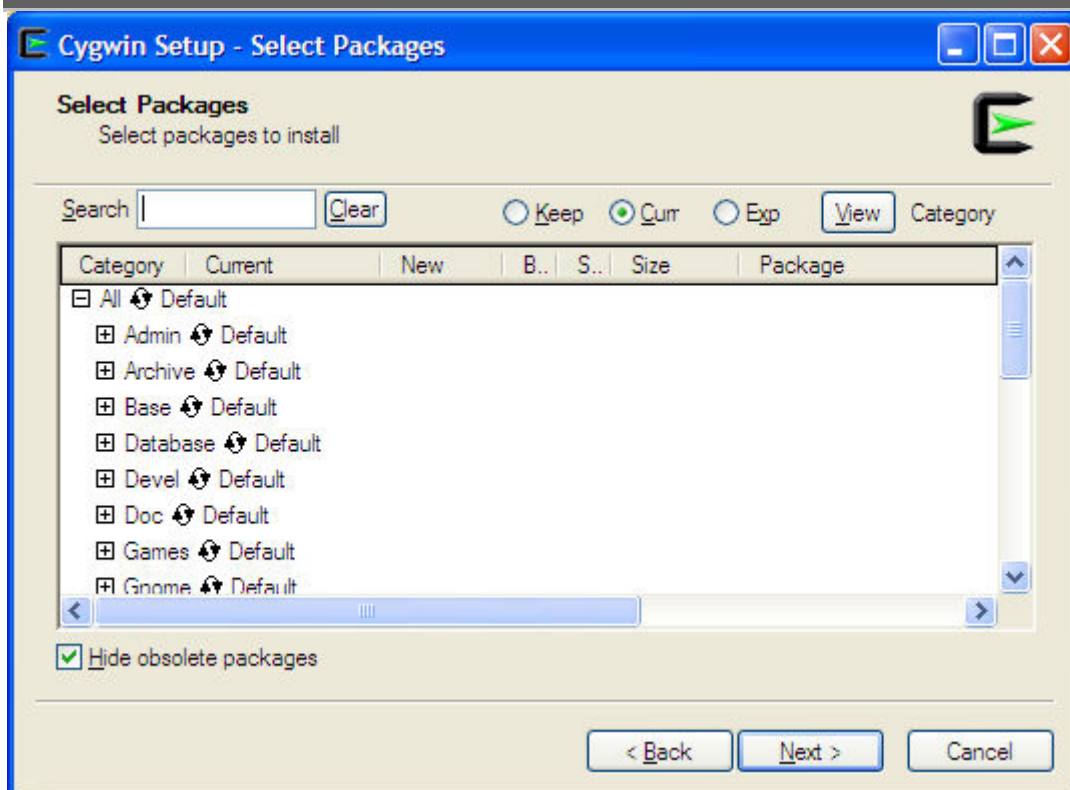
1. Installation & Configurations
2. Show screen shots of installation and configuration on your system which must show your identity with your system.
3. Write a brief summary about installations and configurations on your system.
4. Use simulation files exist in NS 2 library and show also its NAM interface.

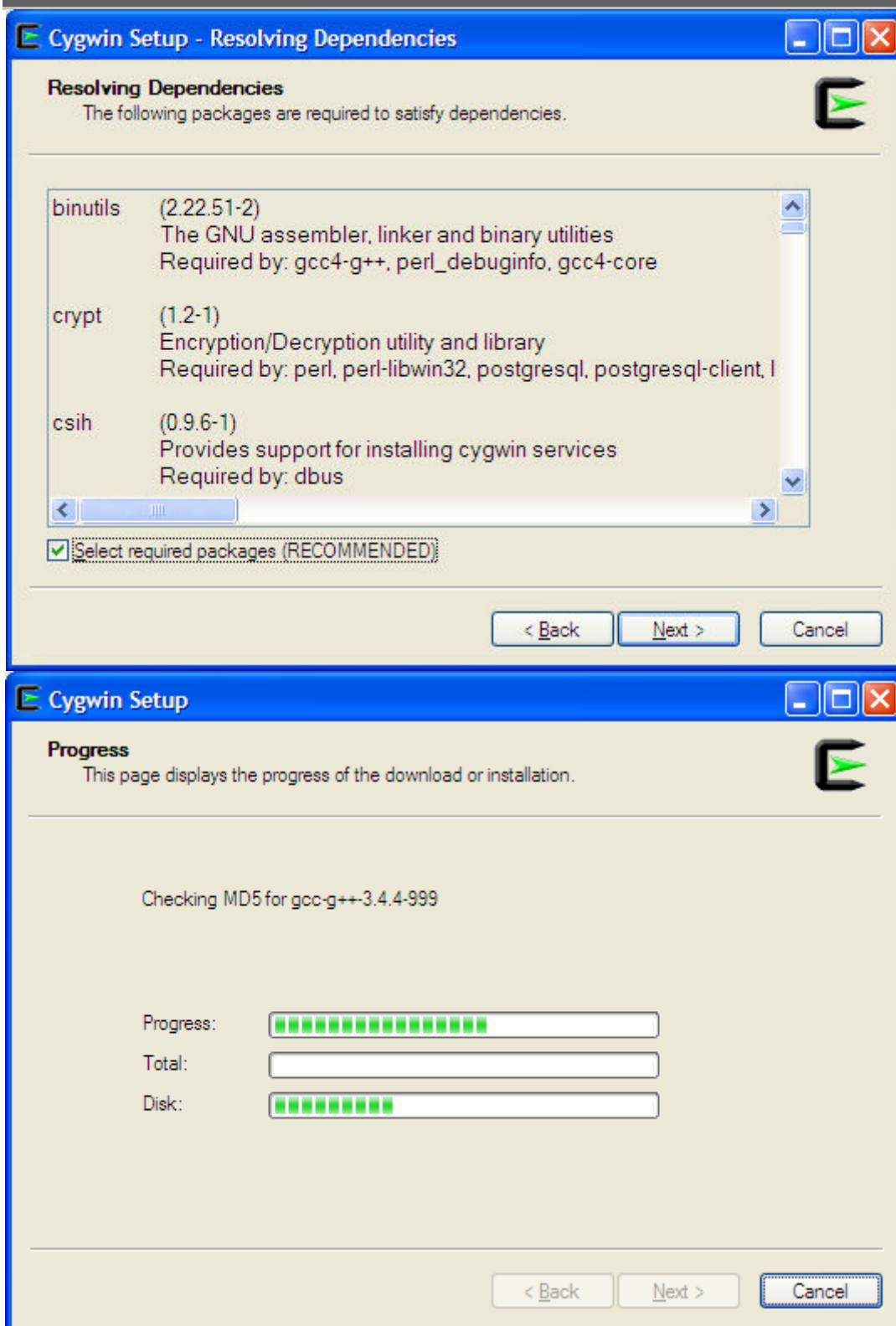
Solution:

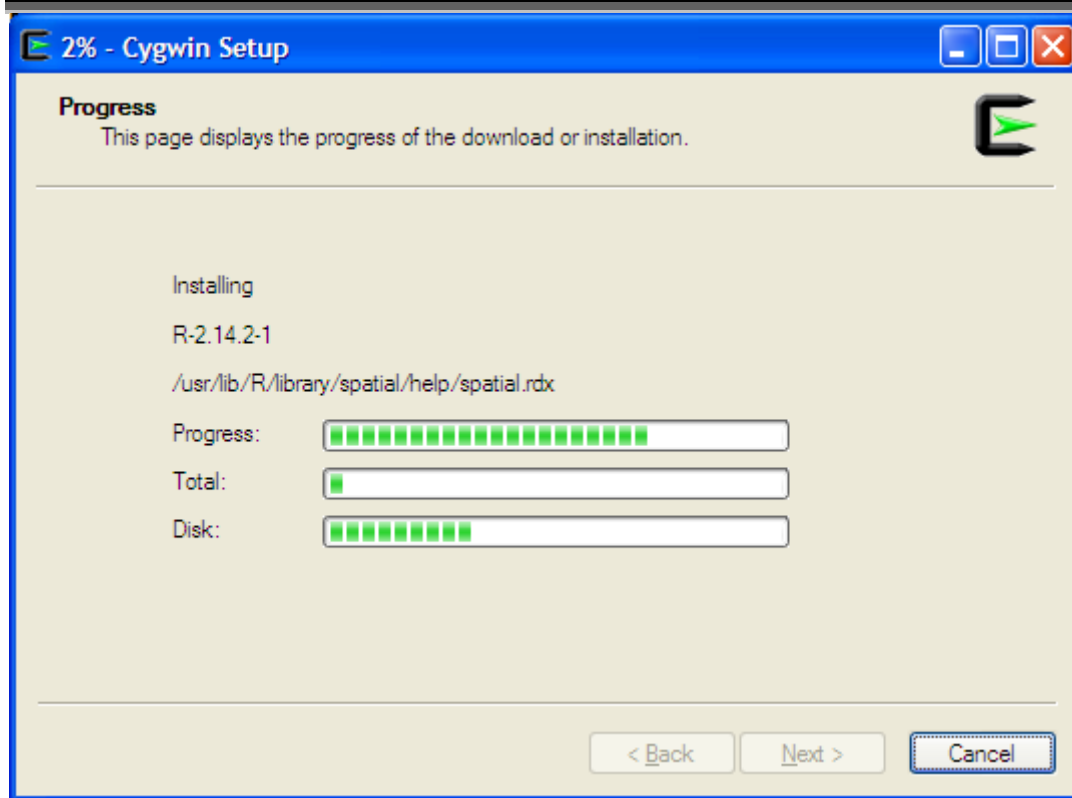
Screenshots of Installation and Configuration

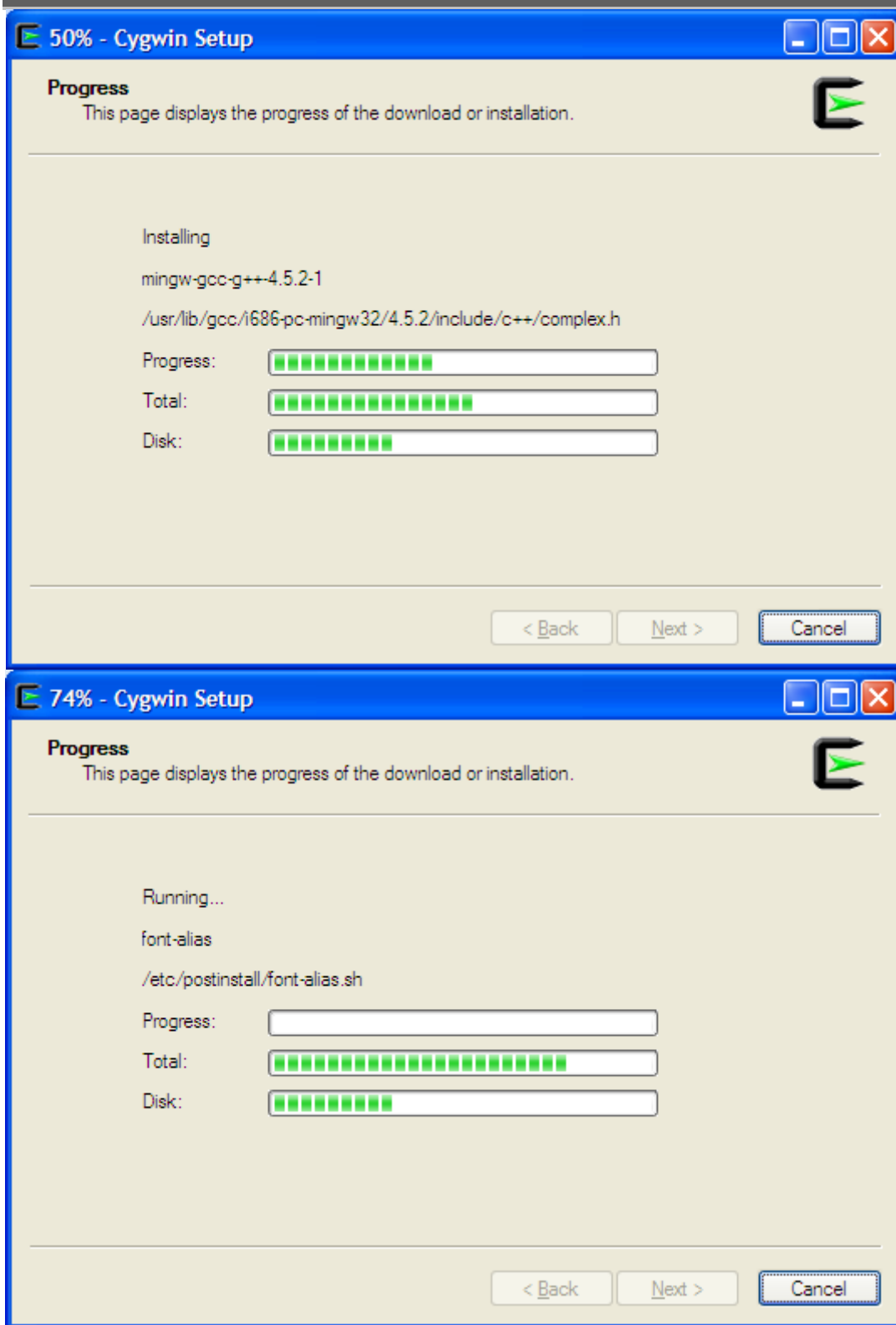


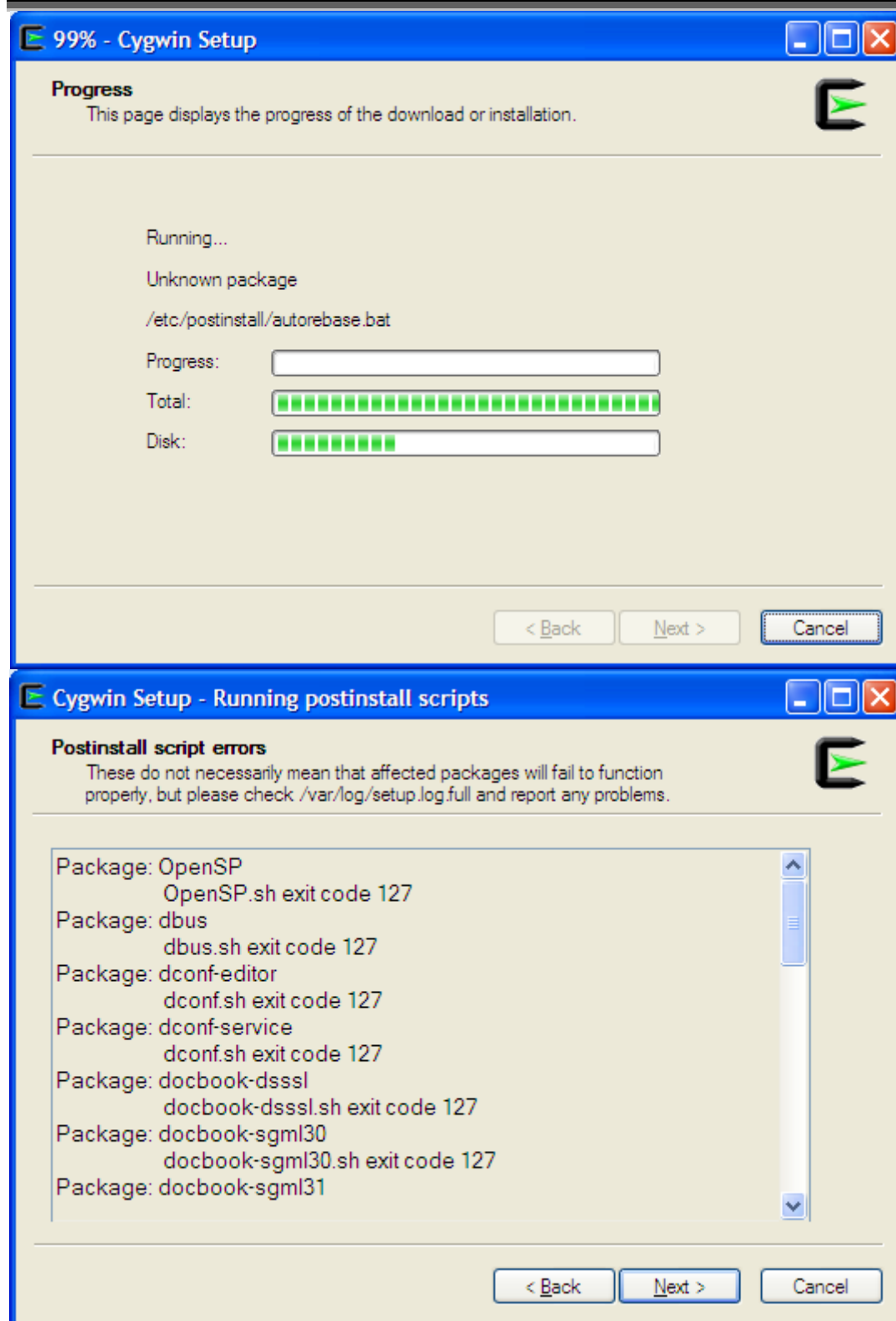


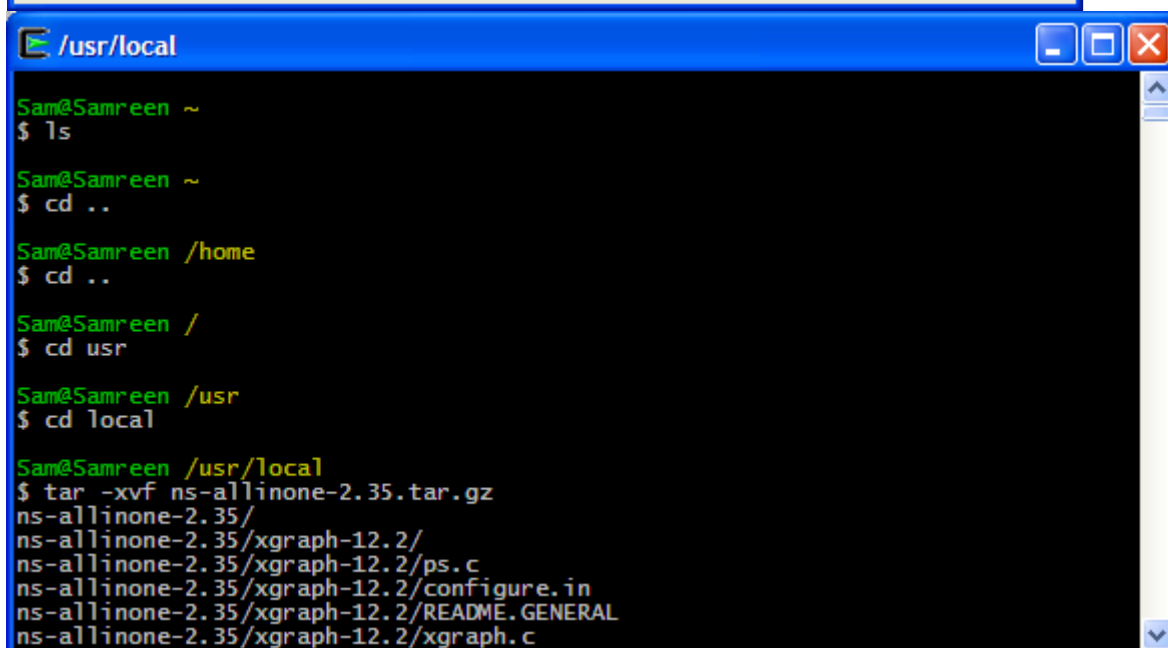
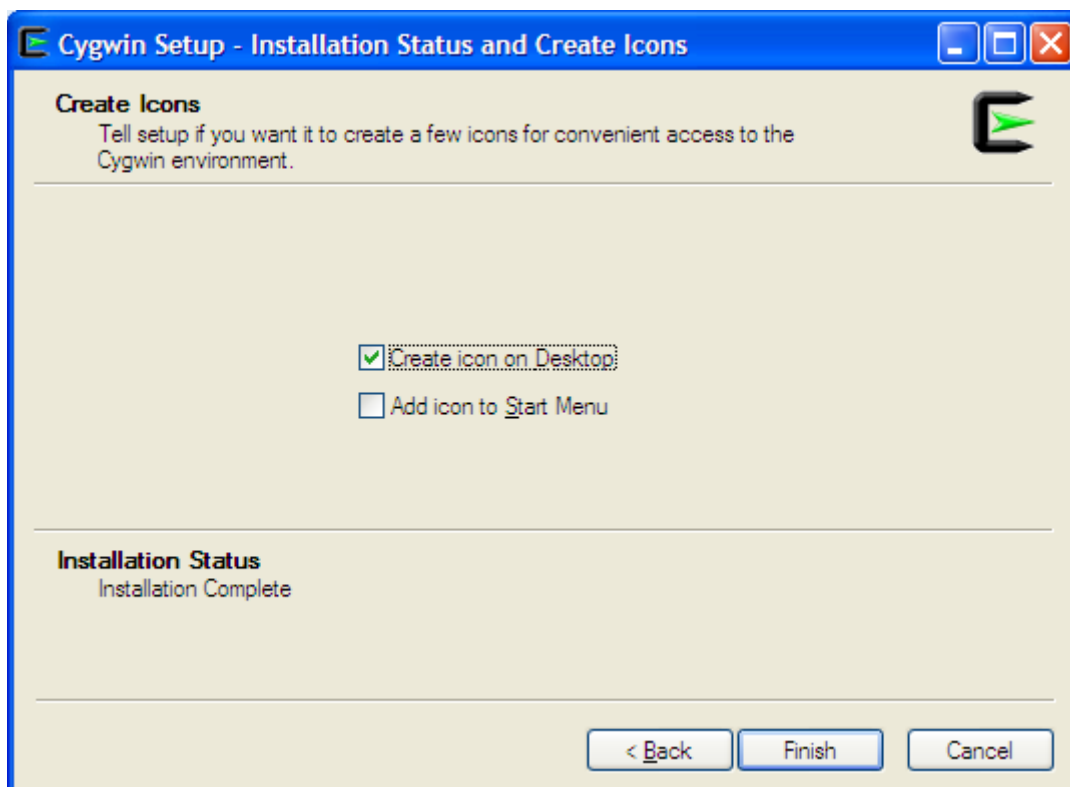












```

/usr/local
ns-allinone-2.35/dei80211mr-1.1.4/src/wireless-channelpa.h
ns-allinone-2.35/dei80211mr-1.1.4/src/phymodes.h
ns-allinone-2.35/dei80211mr-1.1.4/src/libinit.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/peerstatsdb_static.h
ns-allinone-2.35/dei80211mr-1.1.4/src/peerstats.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/test_power_profile.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/rbar.h
ns-allinone-2.35/dei80211mr-1.1.4/src/wireless-phymr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/ra-snr.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-802_11mr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/Mac80211EventHandler.h
ns-allinone-2.35/dei80211mr-1.1.4/src/PER.h
ns-allinone-2.35/dei80211mr-1.1.4/src/power_profile.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.h
ns-allinone-2.35/dei80211mr-1.1.4/src/PER.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/papropagation.h
ns-allinone-2.35/dei80211mr-1.1.4/src/per_table_80211b_intersil_HFA3861B.tcl
ns-allinone-2.35/dei80211mr-1.1.4/src/rateadapter.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/ra-snr.cc
ns-allinone-2.35/install

Sam@Samreen /usr/local
$

/usr/local/ns-allinone-2.35
ns-allinone-2.35/dei80211mr-1.1.4/src/wireless-phymr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/ra-snr.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-802_11mr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/Mac80211EventHandler.h
ns-allinone-2.35/dei80211mr-1.1.4/src/PER.h
ns-allinone-2.35/dei80211mr-1.1.4/src/power_profile.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.h
ns-allinone-2.35/dei80211mr-1.1.4/src/PER.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/papropagation.h
ns-allinone-2.35/dei80211mr-1.1.4/src/per_table_80211b_intersil_HFA3861B.tcl
ns-allinone-2.35/dei80211mr-1.1.4/src/rateadapter.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/ra-snr.cc
ns-allinone-2.35/install

Sam@Samreen /usr/local
$ ls
bin  etc  lib  ns-allinone-2.35  ns-allinone-2.35.tar.gz

Sam@Samreen /usr/local
$ cd ns-allinone-2.35

Sam@Samreen /usr/local/ns-allinone-2.35
$ |

```

```

/usr/local/ns-allinone-2.35
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.h
ns-allinone-2.35/dei80211mr-1.1.4/src/PER.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/papropagation.h
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ns-allinone-2.35/dei80211mr-1.1.4/src/rateadapter.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/ra-snr.cc
ns-allinone-2.35/install

Sam@Samreen /usr/local
$ ls
bin  etc  lib  ns-allinone-2.35  ns-allinone-2.35.tar.gz

Sam@Samreen /usr/local
$ cd ns-allinone-2.35

Sam@Samreen /usr/local/ns-allinone-2.35
$ ls
cweb      install      ns-2.35      sgb          tk8.5.10
dei80211mr-1.1.4  INSTALL.WIN32  otcl-1.14    tcl8.5.10    xgraph-12.2
gt-itm     nam-1.15     README      tclcl-1.20   zlib-1.2.3

Sam@Samreen /usr/local/ns-allinone-2.35
$ |

```

```

/usr/local/ns-allinone-2.35
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.h
ns-allinone-2.35/dei80211mr-1.1.4/src/PER.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/papropagation.h
ns-allinone-2.35/dei80211mr-1.1.4/src/per_table_80211b_intersil_HFA3861B.tcl
ns-allinone-2.35/dei80211mr-1.1.4/src/rateadapter.h
ns-allinone-2.35/dei80211mr-1.1.4/src/mac-timersmr.cc
ns-allinone-2.35/dei80211mr-1.1.4/src/ra-snr.cc
ns-allinone-2.35/install

Sam@Samreen /usr/local
$ ls
bin  etc  lib  ns-allinone-2.35  ns-allinone-2.35.tar.gz

Sam@Samreen /usr/local
$ cd ns-allinone-2.35

Sam@Samreen /usr/local/ns-allinone-2.35
$ ls
cweb      install      ns-2.35      sgb          tk8.5.10
dei80211mr-1.1.4  INSTALL.WIN32  otcl-1.14    tcl8.5.10    xgraph-12.2
gt-itm     nam-1.15     README      tclcl-1.20   zlib-1.2.3

Sam@Samreen /usr/local/ns-allinone-2.35
$ ./install

```

Installation and Configuration

Download zip file of ns2.35 (allinone) from:

<http://sourceforge.net/projects/nsnam/files/allinone/ns-allinone-2.35/>

Download cygwin setup.exe from www.cygwin.com

Cygwin Installation

1. Click on setup.exe of Cygwin
2. A welcome screen will appear, click next
3. A screen will appear asking for the installation type. Select “Install from Local Directory” and click next.
4. A screen will appear asking for the installation directory where Cygwin will be installed. Select Install for “All users” and click next.
5. A screen will appear asking for the selection of local package directory where downloaded Cygwin files are located. Choose the appropriate directory and click next.
6. A screen will appear asking for the packages to install. Initially the “Default” option is selected. Change it to “Install” and make sure that the following packages are necessarily selected for installation: gcc4, gcc4-g++, gawk, tar, gzip, make, patch, perl, w32api, xorg-server, xinit, libX11-devel, libXmu-devel. Click next.
7. A screen will appear for “Resolving Dependencies”. Click next.
8. Next, the Installation will start and the window will show the progress. Wait until the installation is finished. This will take approximately 1 hour or so.
9. After, the installation is finished; a screen will appear with the note “installation completed” and asking for the options of creating desktop icons and start menu entries. Select the desired option and click “Finish”.
10. Cygwin is installed now.
11. Open Cygwin through the desktop icon created by the installation.
12. It will show the computer name and the name of person it is registered to in the format “Registered to@Computer name”. In my case it is “Sam@Samreen”.

NS2 Installation

1. Now, change the directory to usr/local by using the cd command.
2. Copy ns-allinone-2.35.tar.gz to directory C:/cygwin/usr/local
3. Extract the ns2 archive into the usr/local folder.
4. Now change the directory to “ns-allinone-2.35” using the cd command.
5. Type ls to see the list of files in the folder.
6. Type ./install and press enter.
7. The NS-2 installation will start. It will start after making sure the required packages mentioned above are installed. Otherwise it will give error and will ask if the user wants to proceed anyway.
8. Next, the Installation will start and the Console window will show the progress. Wait until the installation is finished. This will also take approximately 1 hour or so.
9. Configure system variables
10. Start Cygwin using the startxwin command
11. Write “ns ex1.tcl” to execute the ns-2 nam interface.